

Problem 2. (26 points) Consider the graph G , given by its (symmetric) adjacency matrix:

$$G = \begin{bmatrix} 0 & & & & & \\ 3 & 0 & & & & \\ 6 & 4 & 0 & & & \\ \infty & 2 & 5 & 0 & & \\ 12 & 8 & 3 & 1 & 0 & \end{bmatrix}$$

- a. (1 pts) Fill in the rest of the adjacency matrix.
- b. (1 pts) Is this a planar graph?

- c. (3 pts) Draw the graph in the top right space, using node labels $\{A, B, C, D, E\}$, where A corresponds to the first node of the adjacency matrix, etc. Put the nodes clockwise on a pentagon, with A at the top. If the graph is planar, avoid all arc-intersections; otherwise, minimize their number.
- d. (1 pts) Give an adjacency list representation of node A .

- e. (4 pts) Bellman-Ford can be written so that it “settles” one node at each iteration as paths of the next longer length are added. The first node settled is the start node, at paths of length 0. We can then eliminate a settled node from further consideration, including all references to the “settled” node from the adjacency lists of the remaining nodes, effectively reducing the number of nodes in the graph by 1. Here’s a pseudo-code algorithm:
 - i. Let Bellman-Ford choose node x to eliminate; then
 - ii. For each node y in the adjacency list of x , do:
search for and destroy the arc referencing x in y ’s adjacency list.

Suppose that the algorithm we use for “search and destroy” uses a sequential search as it peruses an adjacency list. Considering a simple graph of n nodes, what is the worst case number of comparisons the algorithm might have to do in carrying out Bellman-Ford completely?

Problem 2, cont.

- f. (6 pts) Perform Dykstra's algorithm on the graph G , starting at the node A and going to the node E . Give the initial values of d and s , and then their values after the first two iterations (if necessary – note: the adjacency matrix appears below, with Floyd's):

	A	B	C	D	E
d					
s					

	A	B	C	D	E
d					
s					

	A	B	C	D	E
d					
s					

- g. (6 pts) Use Bellman-Ford starting from node A , giving the initial values of d and s and then their values after the first two iterations:

	A	B	C	D	E
d					
s					

	A	B	C	D	E
d					
s					

	A	B	C	D	E
d					
s					

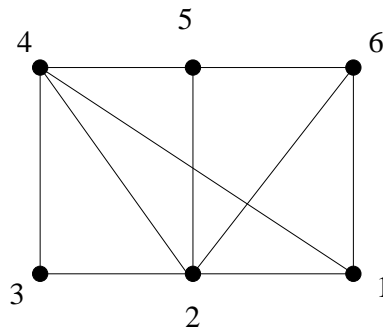
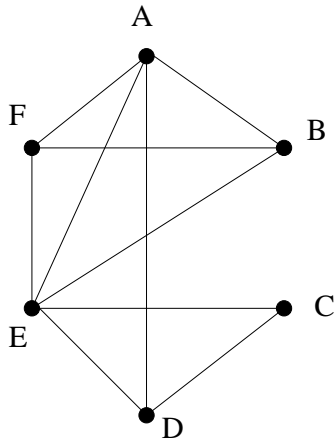
- h. (4 pts) Use Floyd's to give the updated adjacency matrix after finishing the $k = 1$ case (checking for shortcuts via A):

Initial:
$$\begin{bmatrix} 0 & & & & & \\ 3 & 0 & & & & \\ 6 & 4 & 0 & & & \\ \infty & 2 & 5 & 0 & & \\ 12 & 8 & 3 & 1 & 0 & \end{bmatrix}$$

After $k=1$:
$$\begin{bmatrix} 0 & & & & & \\ & 0 & & & & \\ & & 0 & & & \\ & & & 0 & & \\ & & & & 0 & \\ & & & & & 0 \end{bmatrix}$$

Problem 3. (14 points total)

- a. (6 pts): Demonstrate whether the two graphs are isomorphic or non-isomorphic



- b. Show all possible distinct ways to inter-connect identical computers (with single cables – not connecting two is an acceptable choice, too). There is one way to connect one computer: it is connected to nothing else. So we don't connect computers to themselves.

i. (2pts) two computers:

ii. (2pts) three computers:

iii. (4pts) four computers:

- c. **Extra credit** (2pts): Can you figure out a rule for the number of isomorphic simple graphs with n nodes?

Problem 5. (10 points) True or False? (Fix any false statements.)

- a. () A Hamiltonian Circuit exists for any complete graph with at least 3 nodes.
- b. () If there is an Euler path on a graph, the graph can be drawn without lifting one's pencil from the paper.
- c. () The power set of the empty set is empty.
- d. () The integers and the real numbers are the same size as sets.
- e. () Binary search is optimal in its worst case behavior.