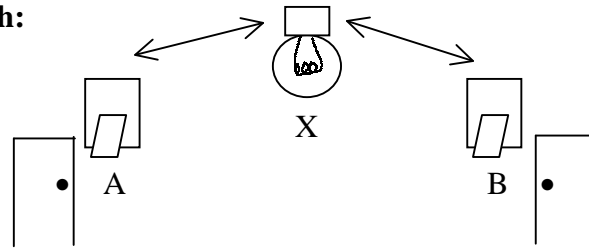


• **3-way light switch:**



- want either switch to be able to turn light ON/OFF independent of position of other switch

$X = \begin{cases} 1, \text{ light ON} \\ 0, \text{ OFF} \end{cases}$, $A = \begin{cases} 1, \text{ switch UP} \\ 0, \text{ switch DOWN} \end{cases}$, $B = \begin{cases} 1, \text{ UP} \\ 0, \text{ DOWN} \end{cases}$

- assume when first installed $X = 0, A = 0,$ and $B = 0$

- want: $A = 0$ and $B = 0 \Rightarrow X = 0$

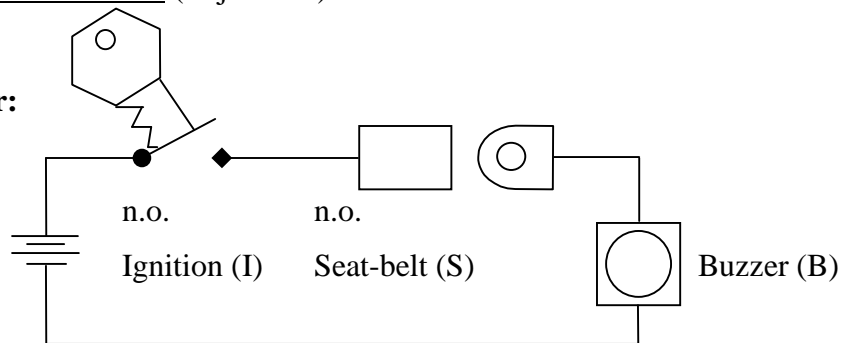
$A = 0 \rightarrow 1$ and $B = 0 \Rightarrow X = 1$	} \Rightarrow	A	B	X	} \Rightarrow circuit (6)
$A = 1$ and $B = 0 \rightarrow 1 \Rightarrow X = 0$		0	0	0	
$A = 1 \rightarrow 0$ and $B = 1 \Rightarrow X = 1$		0	1	1	
$A = 0$ and $B = 1 \rightarrow 0 \Rightarrow X = 0$		1	0	1	
		1	1	0	

- used double pole switches to realize (6) : $X = \bar{A} \cdot B + A \cdot \bar{B}$

- (6) termed exclusive-OR (XOR) since A or $B = 1$ but not both

- (7) termed inclusive-OR (or just OR) since A or B or both = 1

• **Seat-belt buzzer:**



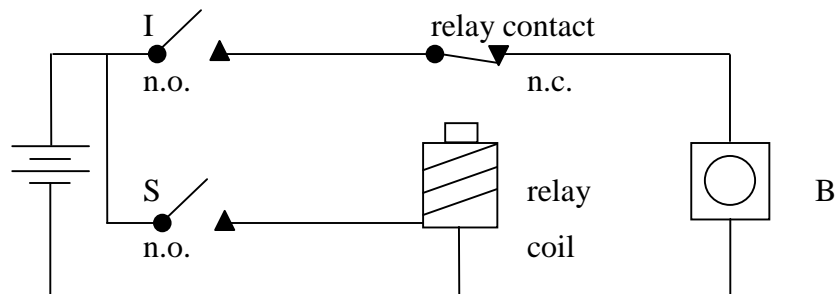
- want:

$I = \begin{cases} 1, \text{ ignition ON} \\ 0, \text{ OFF} \end{cases}$, $S = \begin{cases} 1, \text{ fastened} \\ 0, \text{ unfastened} \end{cases}$, $B = \begin{cases} 1, \text{ sound} \\ 0, \text{ quiet} \end{cases}$

I	S	B
0	0	0
0	1	0
1	0	1
1	1	0

\Rightarrow circuit ② $\Rightarrow B = I \cdot \bar{S} \Rightarrow$ seat-belt needs to be n.c. switch

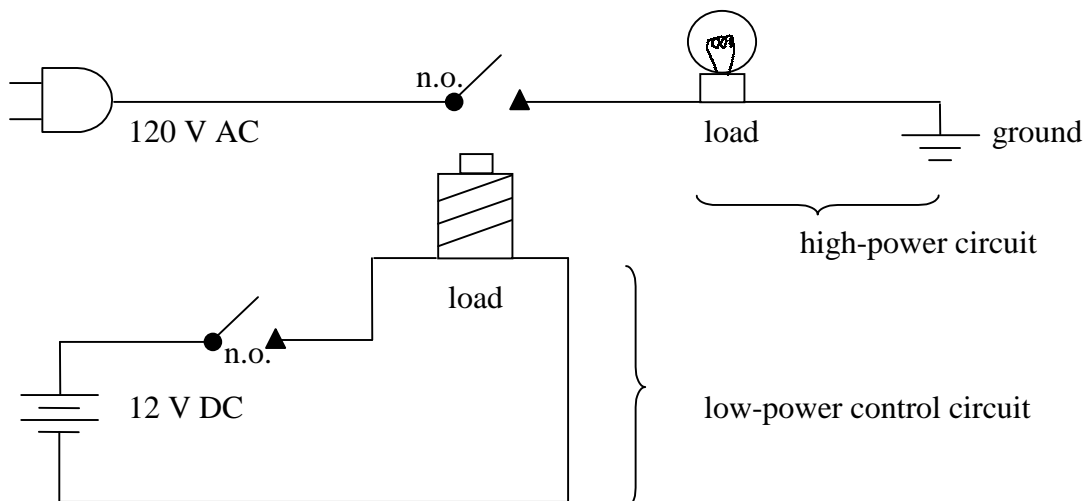
- physically, S best as n.o., but want it to be n.c. in control circuit
- \Rightarrow can use relay:



- both B and relay coil are loads
- when S actuated, current flows through relay coil, which becomes a magnet that actuates relay contact

• **Relays:**

- up to now, all inputs have been mechanical switches that directly made or broke the circuit controlling the load
- relays can be used to indirectly control higher power circuit:



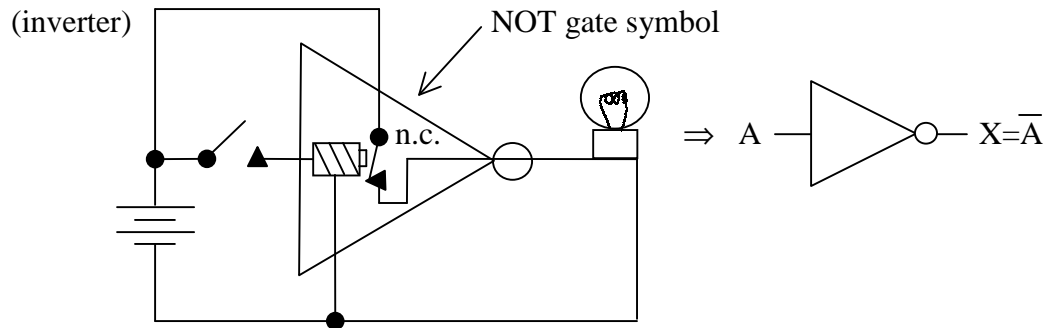
- relays can be normally open (n.o.) or normally closed (n.c.)
- use of relays eliminates need for complex multiple pole switches and wiring

- **Logic Gates:** (see handout)

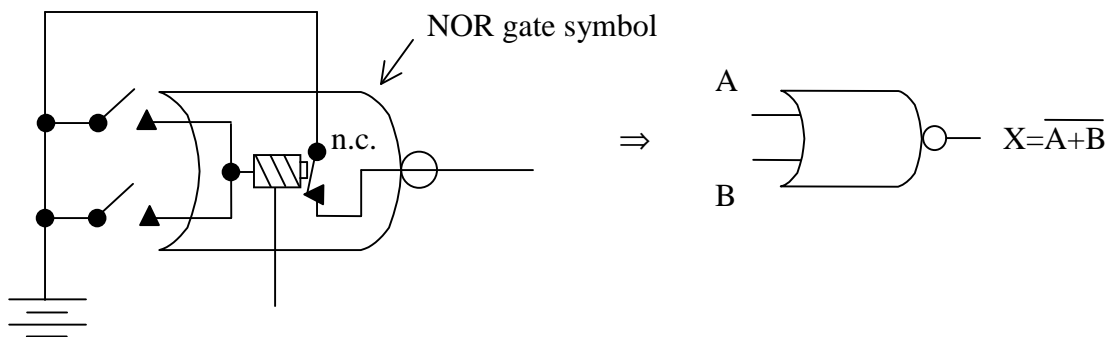
- in seat-belt control circuit, relay used to convert n.o. mechanical switch to n.c. switch

⇒ performing Logical NOT ⇒ inverting signal ⇒ inverter

- **NOT gate:**



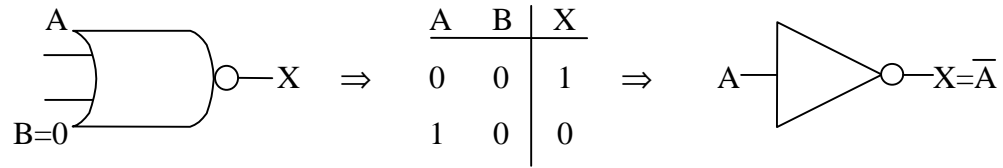
- **NOR gate:** when two (or more) inputs connected in parallel to single relay of NOT gate, get NOR gate



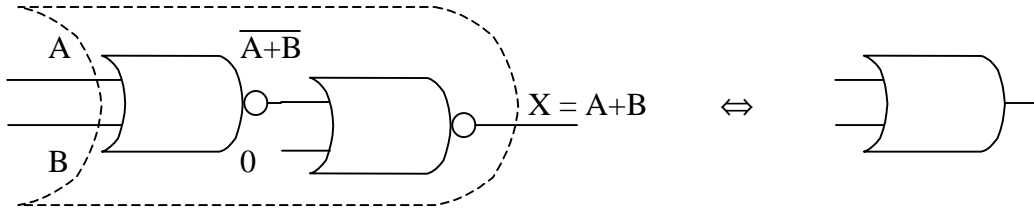
- if n.o. relay contact used in NOR gate ⇒ **OR gate** = A, B — OR — X = A+B

- but, unlike OR gate, NOR gate (and NAND gate) are universal gates
- ⇒ any Boolean logic circuit can be realized using just NOR gates (or just NAND gates)

- NOT gate via NOR gate:



- OR gate via NOR gates:



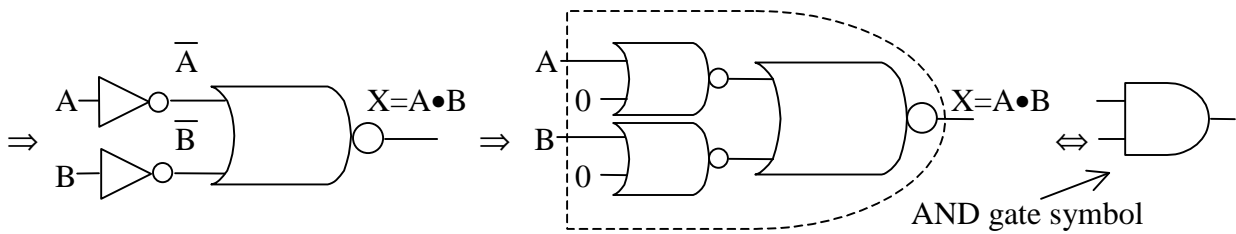
- AND gate via NOR gates: not immediately obvious how to construct
 \Rightarrow use Theorems of Boolean Algebra

- **Theorems of Boolean Algebra** (see handout)

- used to manipulate Boolean expressions
- developed by George Boole in 1850s (Laws of Thought)
- in 1938, Claude Shannon saw one-to-one correspondence between Boolean expressions and switching circuits
 \Rightarrow Boolean algebra can be used to simplify logic control circuits
- AND gate via NOR gates + Boolean algebra:

Want: $X = A \bullet B$
 $= \overline{\overline{A \bullet B}}$ (by Negation Th.)

Let $A' = \overline{A}$ and $B' = \overline{B}$
 $X = \overline{\overline{A' \bullet B'}}$
 $= \overline{A' + B'}$ (by DeMorgan 2 — left side of DeMorgan 2 is NOR)
 $= \overline{\overline{A} + \overline{B}}$
 $= A + B$



- Theorems also give NOT and OR gates via NOR gates:

NOT: $X = \bar{A}$

$= \bar{A} \bullet 1$ (by Char. 2)

$= \bar{A} \bullet \bar{0}$ (1=0 & 0=1)

$= \overline{A + 0}$ (by DeMorgan 2)

OR: $X = A + B$

$= \overline{\overline{A} + \overline{B}}$ (by Neg.)

$= \overline{\overline{A} \bullet \overline{B}}$ (by DeMorgan 1)

$= \overline{\overline{A} \bullet \overline{B} \bullet 1}$ (by Char. 2)

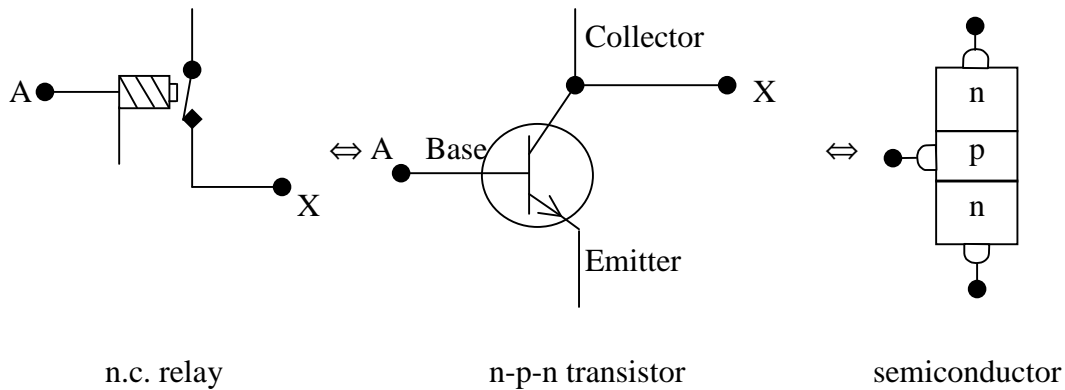
$= \overline{\overline{A} \bullet \overline{B} \bullet \bar{0}}$ (1=0)

$= \overline{\overline{A} \bullet \overline{B} + 0}$ (by DeMorgan 2)

$= \overline{\overline{A + B} + 0}$ (by DeMorgan 2)

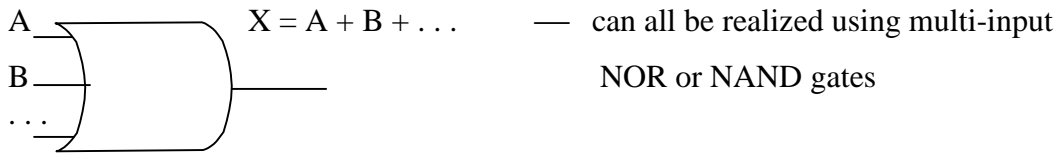
- **Transistor Logic:**

- Why are 3 NOR gates used instead of 1 AND gate?
- NOR gate easy to make using transistor
- transistors used instead of relays in all control applications except to switch high power circuits

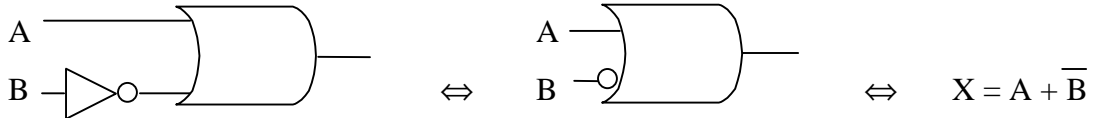


- for control system design, can think in terms of relays

- **Multi-input gates:** logic gates can have more than two inputs

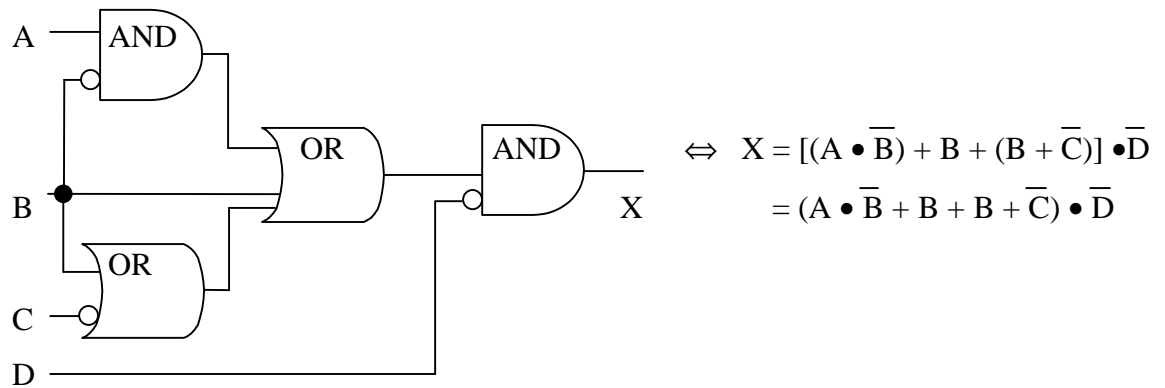
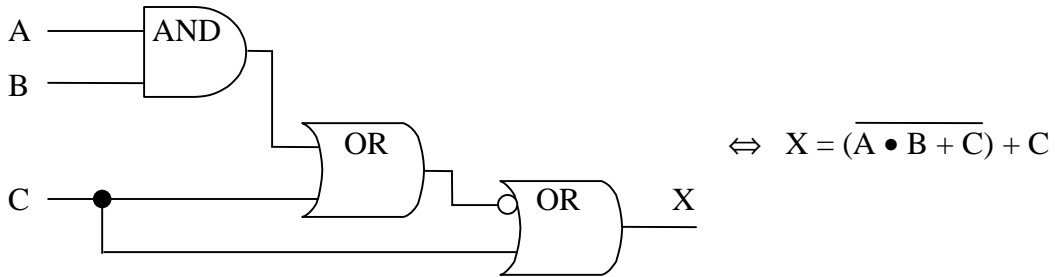


- **Input negation:**



- **Nesting:** when output from gate used as input to next, equivalent to parenthesis in Boolean expression (AND higher precedence than OR)

- **Logic gate network** \Rightarrow Boolean expression:



Theorems and Laws of Boolean Algebra

CHARACTERISTIC THEOREMS

1. $X \cdot 0 = 0$
2. $X \cdot 1 = X$
3. $X + 0 = X$
4. $X + 1 = 1$

COMMUTATIVE LAW

1. $X + Y = Y + X$
2. $X \cdot Y = Y \cdot X$

ASSOCIATIVE LAW

1. $X + Y + Z = X + (Y + Z)$
 $= (X + Y) + Z$
2. $X \cdot Y \cdot Z = X \cdot (Y \cdot Z)$
 $= (X \cdot Y) \cdot Z$

DISTRIBUTIVE LAW

1. $X \cdot Y + X \cdot Z = X(Y + Z)$
2. $(X + Y)(W + Z) = XW + XZ + YW + YZ$

IDEMPOTENT THEOREMS

1. $X \cdot X = X$
2. $X + X = X$

NEGATION THEOREM

$$\overline{\overline{X}} = X$$

INCLUSION THEOREMS

1. $X \cdot \overline{X} = 0$
2. $X + \overline{X} = 1$

ABSORPTIVE LAWS

1. $X + XY = X$
2. $X(X + Y) = X$

REFLECTIVE THEOREMS

1. $X + \overline{X}Y = X + Y$
2. $X(\overline{X} + Y) = XY$
3. $XY + \overline{X}YZ = XY + YZ$

CONSISTENCY THEOREM

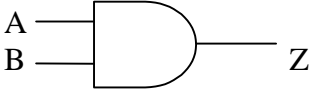
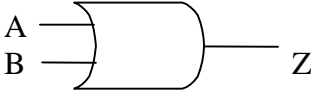
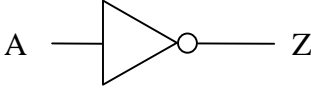
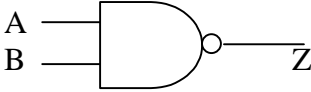
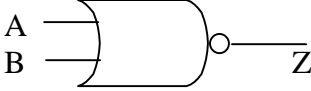
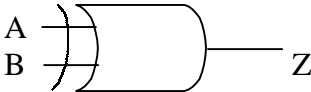
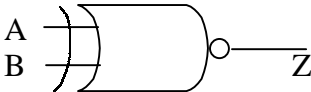
1. $XY + X\overline{Y} = X$
2. $(X + Y)(X + \overline{Y}) = X$

DEMORGAN'S LAWS

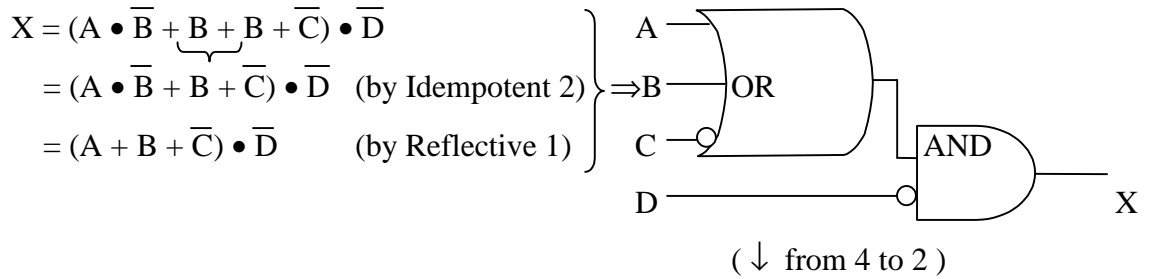
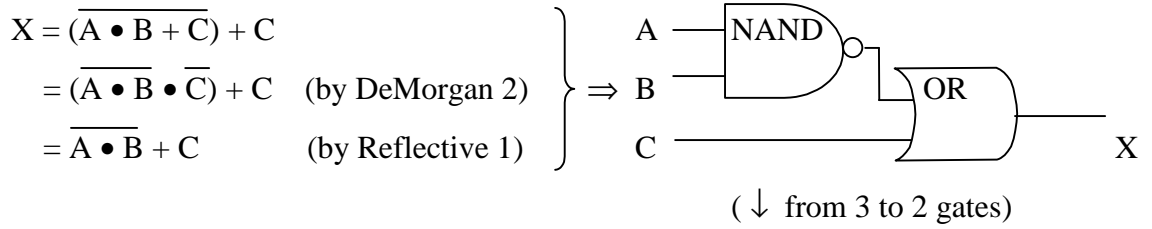
1. $\overline{XY} = \overline{X} + \overline{Y}$
 2. $\overline{X + Y} = \overline{X}\overline{Y}$
-

Logic:	\wedge	A	B	\vee	\downarrow	\leftrightarrow	\equiv	\bar{B}	\leftarrow	\bar{A}	\supset	\rightarrow	\uparrow				
Boolean:	\bullet			\oplus	$+$	\odot											
Name:	AND			XOR	OR	NOR	XNOR	NOT		NOT		NAND					
A	B	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
	short circuit	$A \bullet B$	$A \bullet \bar{B}$	A	$\bar{A} \bullet B$	B	$A \bullet \bar{B} + \bar{A} \bullet B$	$A+B$	$\overline{A \bullet B}, \overline{A+B}$	$A \bullet B + \bar{A} \bullet \bar{B}$	\bar{B}	$A+\bar{B}$	\bar{A}	$\bar{A}+B$	$\overline{\bar{A}+B}, \overline{A \bullet B}$	no control	

**DIGITAL LOGIC GATES AND ASSOCIATED LOGICAL OPERATIONS FOR
BINARY VARIABLES**

Name	Symbol	Logical operation	Truth table															
AND		$Z = A \cdot B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	Z	0	0	0	0	1	0	1	0	0	1	1	1
A	B	Z																
0	0	0																
0	1	0																
1	0	0																
1	1	1																
OR		$Z = A + B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	Z	0	0	0	0	1	1	1	0	1	1	1	1
A	B	Z																
0	0	0																
0	1	1																
1	0	1																
1	1	1																
NOT		$Z = \bar{A}$	<table border="1"> <thead> <tr> <th>A</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	Z	0	1	1	0									
A	Z																	
0	1																	
1	0																	
NAND		$Z = \overline{A \cdot B}$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	Z	0	0	1	0	1	1	1	0	1	1	1	0
A	B	Z																
0	0	1																
0	1	1																
1	0	1																
1	1	0																
NOR		$Z = \overline{A + B}$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	Z	0	0	1	0	1	0	1	0	0	1	1	0
A	B	Z																
0	0	1																
0	1	0																
1	0	0																
1	1	0																
XOR		$Z = A \oplus B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	Z	0	0	0	0	1	1	1	0	1	1	1	0
A	B	Z																
0	0	0																
0	1	1																
1	0	1																
1	1	0																
XNOR		$Z = A \odot B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>Z</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	Z	0	0	1	0	1	0	1	0	0	1	1	1
A	B	Z																
0	0	1																
0	1	0																
1	0	0																
1	1	1																

- Simplifying Boolean expression \Leftrightarrow reducing number of logic gates



- **Canonical Sum-of-Products Form**

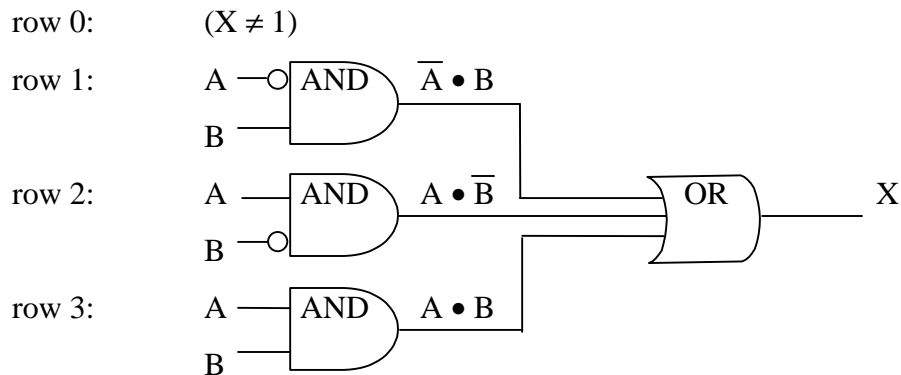
- Used to realize any logic control circuit from its truth table

1. For each circuit, construct its truth table to relate all possible inputs to desired output of circuit:

Row	A	B	X
0	0	0	0
1	0	1	1
2	1	0	1
3	1	1	1

} desired output

2. For each row where $X = 1$, AND together inputs, where 0 inputs are NOTed



3. OR together AND gates from step 2

$$X = \overline{A} \bullet B + A \bullet \overline{B} + A \bullet B$$

4. Use Theorems of Boolean Algebra to simplify if possible

$$= \overline{A} \bullet B + A \bullet (\overline{B} + B) \quad (\text{by Distrib. 1})$$

$$= \overline{A} \bullet B + A \bullet (1) \quad (\text{by Inclusion 2})$$

$$= \overline{A} \bullet B + A \quad (\text{by Char. 2})$$

$$= B + A \quad (\text{by Reflect. 1})$$

$$= A + B \Rightarrow \text{OR} \quad (\text{by Commut. 1})$$

- termed "Sum-of-Products" since OR-ing together (summing) AND-ed inputs (products)

- Example: XOR

A	B	X	}	⇒	X = $\overline{A} \bullet B + A \bullet \overline{B}$, can't simplify
0	0	0			
0	1	1			
1	0	1			
1	1	0			

- Example: Seat belt and Door-Open Buzzer

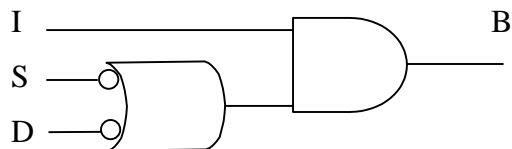
$$I = \begin{cases} 1, \text{ ignition ON} \\ 0, \text{ OFF} \end{cases}, \quad S = \begin{cases} 1, \text{ seat belt fastened} \\ 0, \text{ unfastened} \end{cases}$$

$$D = \begin{cases} 1, \text{ door closed} \\ 0, \text{ door open} \end{cases}, \quad B = \begin{cases} 1, \text{ sound} \\ 0, \text{ quiet} \end{cases}$$

I	S	D	B
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

quiet when
I = 0

$$\begin{aligned}
 B &= I \cdot \bar{S} \cdot \bar{D} + I \cdot \bar{S} \cdot D + I \cdot S \cdot \bar{D} \\
 &= I \cdot (\bar{S} \cdot \bar{D} + \bar{S} \cdot D + S \cdot \bar{D}) \quad (\text{by Distrib. 1}) \\
 &= I \cdot (\bar{S} + S \cdot \bar{D}) \quad (\text{by Consist. 1}) \\
 &= I \cdot (\bar{S} + \bar{D}) \quad (\text{by Reflect. 1})
 \end{aligned}$$



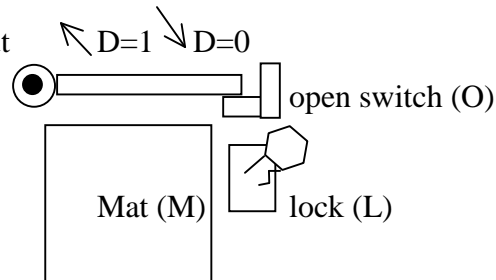
- Example: Automatic Door and Lock

D = 1, door opening
0, door closing

O = 1, door not shut
0, door shut

M = 1, someone on mat
0, no one on mat

L = 1, unlocked
0, locked



Operation: Want door to open if someone on mat and door unlocked.

If locked, want it to stay open or shut independent of mat

	M	L	O	D	
Stay shut	0	0	0	0	start 0
	0	0	1	1	unlock 1
	0	1	0	0	
stay open	0	1	1	0	closing 5
	1	0	0	0	open 2
	1	0	1	1	
	1	1	0	1	opening 3
	1	1	1	1	shut 4

$$\begin{aligned}
 D &= \bar{M} \cdot \bar{L} \cdot O + M \cdot \bar{L} \cdot O + M \cdot L \cdot \bar{O} + M \cdot L \cdot O \\
 &= (\bar{M} + M) \cdot \bar{L} \cdot O + M \cdot L \cdot (\bar{O} + O) \quad (\text{by Distrib. 1}) \\
 &= (1) \cdot \bar{L} \cdot O + M \cdot L \cdot (1) \quad (\text{by Inclus. 2}) \\
 &= \bar{L} \cdot O + M \cdot L
 \end{aligned}$$

