

MAT360 Section Summary: 2.2

Fixed-Point Iteration

Summary

Suppose that you want to solve the equation

$$\cos(x) = x$$

The value of x that satisfies this equation is called a **fixed point** for the function $g(x) = \cos(x)$, because it is a point such that $g(x) = x$ — the image is the same as the argument.

One way to go about finding the fixed point would be to rewrite the equation as

$$f(x) = \cos(x) - x = 0$$

and to use bisection to find a root of f (in fact, the unique root, as one can see from the graph of f).

Fixed-point iteration is based on a couple of results from calculus: the IVT, and the MVT, as follows:

Theorem 2.2:

- If $g \in C[a, b]$ and $g(x) \in [a, b]$ for all $x \in [a, b]$, then g has a fixed point in $[a, b]$.
- If, in addition, $g'(x)$ exists on (a, b) and a positive constant $k < 1$ exists with

$$|g'(x)| \leq k, \quad \forall x \in (a, b),$$

then the fixed point in $[a, b]$ is unique.

The proofs are by

- the IVT, with $h(x) = g(x) - x$; and

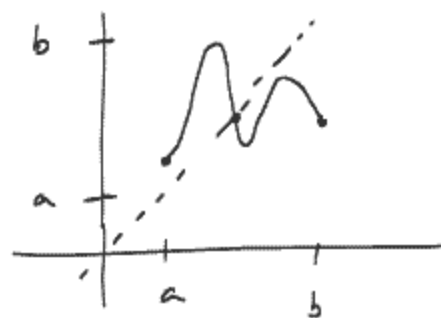
If $g(a) = a$ or $g(b) = b$
then we've found our fixed point!

Otherwise,

$h(a) > 0$ and $h(b) < 0$; h is continuous on $[a, b]$, hence by the IVT, h has a root p .

So $h(p) = 0 \Leftrightarrow g(p) - p = 0 \Leftrightarrow \boxed{g(p) = p}$

— p is a fixed point of g .



- the MVT, and contradiction.

Assume that there are two ^{distinct} fixed points,
 $p_1 \neq p_2$ - $g(p_1) = p_1$ + $g(p_2) = p_2$

$$\underbrace{|g(p_1) - g(p_2)|}_{|p_1 - p_2|} = |g'(c)(p_1 - p_2)| \quad \begin{array}{l} \text{by the} \\ \text{MVT} \end{array}$$

$$|p_1 - p_2| = |g'(c)| |p_1 - p_2| \quad c \in [p_1, p_2]$$

$$\Rightarrow |g'(c)| = 1$$

But by assumption, $|g'(x)| < 1$ for $(-1, 1)$; this
 is a contradiction of distinct fixed points.
 Hence the fixed point is unique.

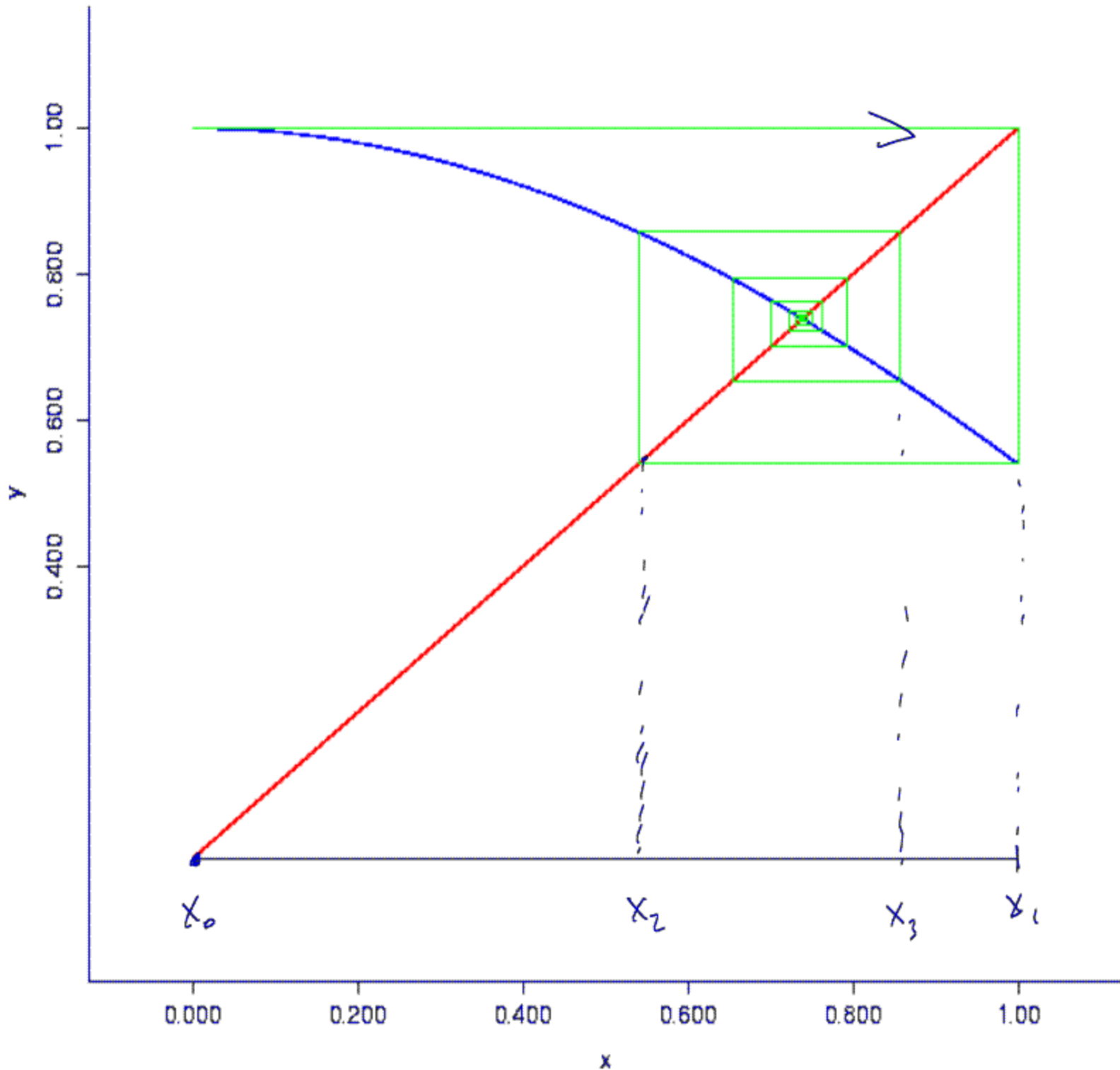
Actually the condition is sufficient but
 not necessary

So we know that there's a fixed point on an interval $[a, b]$, and may even know that it's unique. What now?

Now we assume that, perhaps, if we start with a value x_0 that's close to the real fixed point p , that by simply computing $g(x_0)$ (which is $\approx x_0$) we'll actually get closer to p .

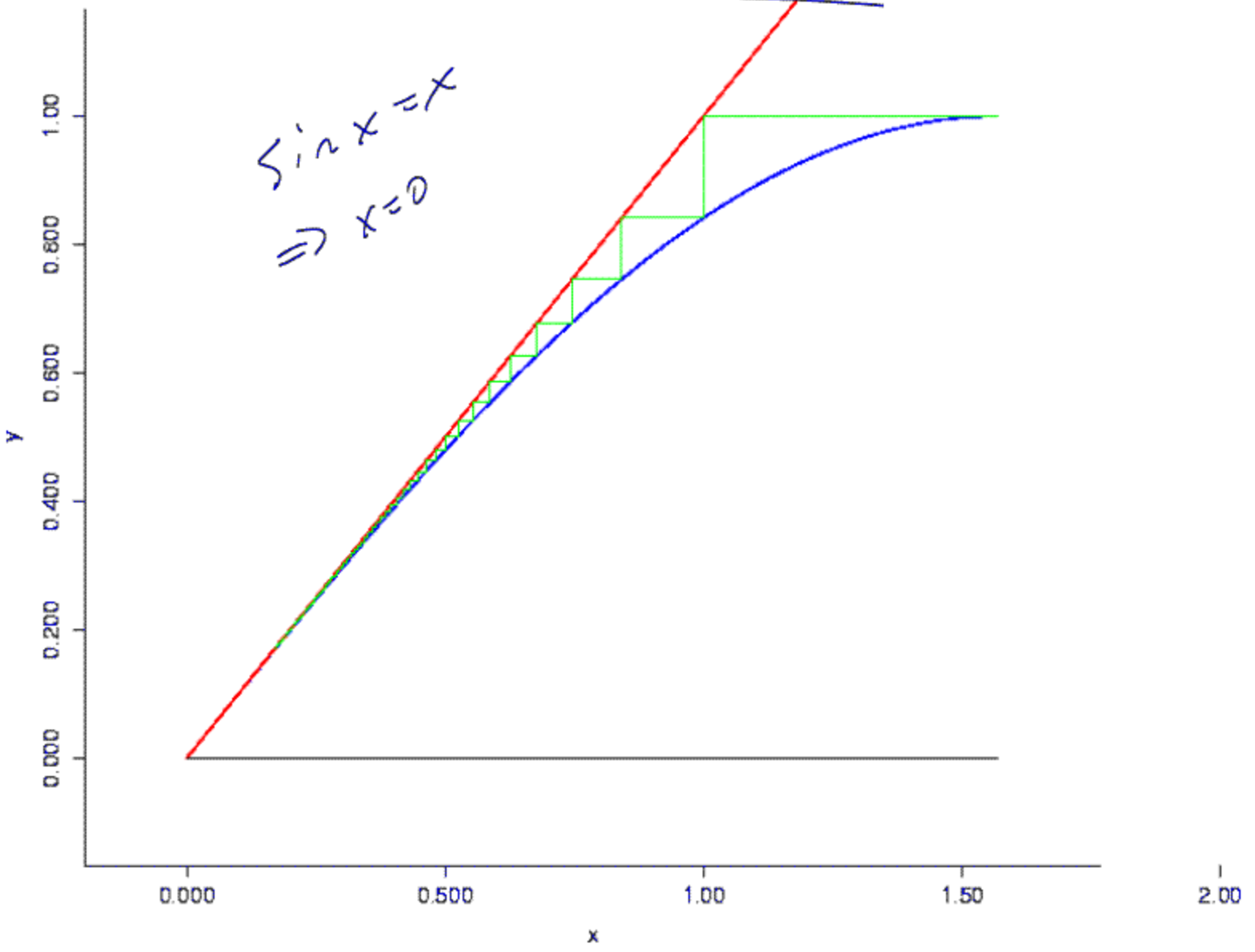
Let's look at the "cobweb diagram" of this situation.

Cobwebbing

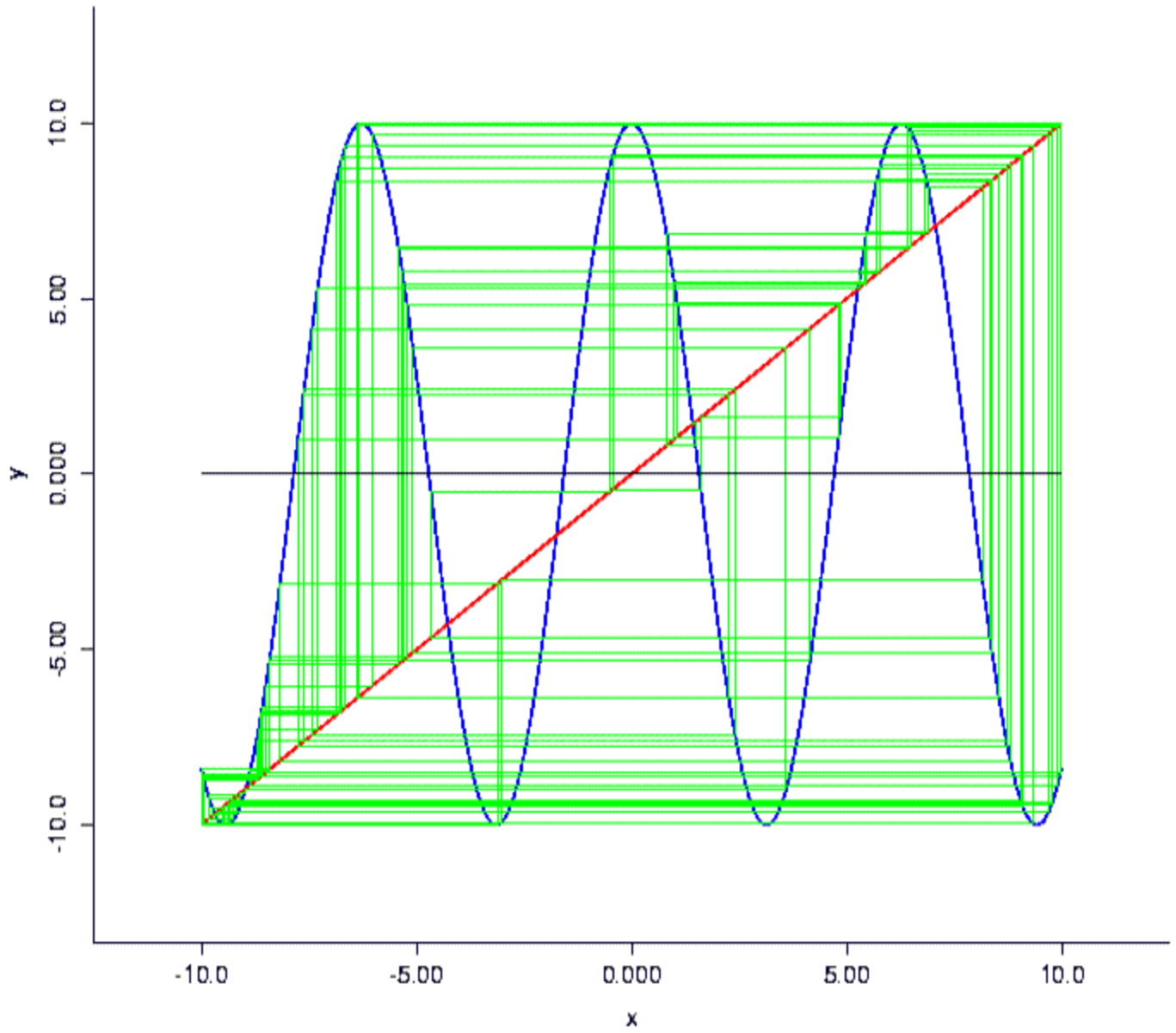


Cobwebbing, 100 iterations max -- fixed point (0.1705, 0.1697)?

$\sin x = x$
 $\Rightarrow x=0$



Cobwebbing, 100 iterations max -- fixed point (9.0973 -9.4687)?



Under what circumstances will that happen? In what circumstances would the same "cobwebbing" procedure fail?

Well, in some circumstances, it's guaranteed to work:

Theorem 2.3: Fixed-Point Theorem Let $g \in C[a, b]$ be such that $g(x) \in [a, b]$, for all x in $[a, b]$. Suppose, in addition, that g' exists on (a, b) , and that a constant $0 < k < 1$ exists with

$$|g'(x)| \leq k$$

for all $x \in (a, b)$. Then for any number $p_0 \in [a, b]$, the sequence

$$p_n = g(p_{n-1})$$

$n \geq 1$, converges to the unique fixed point p in $[a, b]$.

Proof: MVT applied to $|p_n - p|$.

$$|p_n - p| = |g(p_{n-1}) - g(p)| = |g'(c)(p_{n-1} - p)| \leq k |p_{n-1} - p|$$

where c is between p_{n-1} & p

$$|p_n - p| \leq k^n |p_0 - p|$$

Corollary 2.4: If g satisfies the hypotheses of Theorem 2.3, then bounds for the error involved in using p_n to approximate p are given by

$$|p_n - p| \leq k^n \max\{p_0 - a, b - p_0\}$$

and

$$|p_n - p| \leq \frac{k^n}{1-k} |p_1 - p_0| \quad |p_m - p_n| =$$

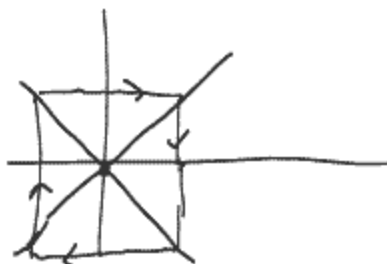
for all $n \geq 1$.

Proof: by use of various inequalities.

There may be lots of ways to create a fixed-point function, and some of them are better than others.

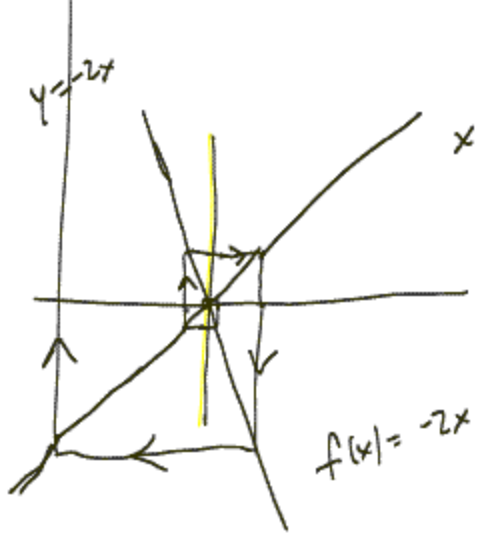
Example: consider exercises 1 and 2.

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$x = x$

"neutrally stable" -
not getting any better,
but not getting
any worse!



Spiralling
out when
 $|\text{slope}| > 1$

#1

$$f(x) = x^4 + 2x^2 - x - 3$$

To solve: $f(x) = 0$

$$g_0(x) = x^4 + 2x^2 - 3 = x$$

$$[g_1(x)]^4 = 3 + x - 2x^2$$

$$x^4 = 3 - x - 2x^2 \Rightarrow f(x) = 0$$

So the solution to the fixed point problem
solves the root problem.