

Overview: Chapters 1-4

Chapter 1: Logic

Propositional: truth tables

Predicate: introduces variables

Useful in defining sets

in calculus $\forall x \exists x$

[Lewis Carroll's Game of Logic
Venn diagrams

Started proving Rings, based
on some properties of $\wedge + \vee$ -

commutativity

associativity

distributivity

complements

identities

(negation)

($T + F, 1 + 0$)

Prove Theorems & put them in our
toolbox.

Two binary
 $\wedge + \vee$
One unary
negation

Chapter 2: proofs & induction

Proof techniques

- direct proof
- by contradiction (Tautology test)
- exhaustion (Four-color problem)
- "not" proof - counterexample
- contraposition

$$P \rightarrow Q \Leftrightarrow Q' \rightarrow P'$$

(distinguish the
contraposition
from the converse).

Induction

First glimpse of infinity
(Natural numbers \mathbb{N})

Two forms of induction:

1st principle - } equivalent
2nd principle - }

Could be finite
Backwards induction
one couldn't be
extended
All horses are the
same color proof
Dominoes

1st principle -

① Base case ($P(1)$)

② $P(n) \rightarrow P(n+1)$

an example of the temporary hypothesis.

Inductive step

if the n^{th} domino falls, then the

$(n+1)^{\text{th}}$ domino falls

$\therefore P(n) \forall n \in \mathbb{N}$ (all dominoes fall)

2nd principle -

① Base case ($P(1)$)

② $P(k) \forall k \mid 1 \leq k \leq n$

$\rightarrow P(n+1)$

All the dominoes have fallen up to n

the $(n+1)^{\text{th}}$ domino falls

$\therefore P(n) \forall n \mid n \in \mathbb{N}$.

Induction concerns only smallest infinity - a countable, denumerable infinity, \mathbb{N} .

Chapter 3 : Recursion

do it; then do it again, & again, & ...
until satisfied.

Idea behind Newton's method
from calculus for solving
 $f(x) = 0$

Defining objects (e.g. wffs).

We've defined trees recursively,
palindromic strings, & even
factorials:

① $0! = 1$ base case

② $P(n+1) = (n+1)P(n)$ inductive step

Definition of $n!$
 $\forall n \in \mathbb{Z}^+$

Recurrence relation

1st order, non-constant
homogeneous.

Another really important one:

Fibonacci numbers.

① $\begin{cases} F(1) = 1 \\ F(2) = 1 \end{cases}$ base cases

② $F(n) = F(n-1) + F(n-2)$ inductive step

Recursion - though a great way to define things - may be a horrible way to compute them!

Analysis of algorithms

- solve some in general
- proved the solution using induction.

(divide & conquer such as binary search).

Euclidean algorithm for finding GCD - Fibonacci #s were the worst case scenario.

Chapter 4: Sets

Three really cool ideas:

① Power sets - the set of all subsets (powers of 2 - Pascal's triangle)

② Infinity & beyond!

There are infinitely many different sizes of infinity - all uncountable after the smallest infinity - \aleph_1 .

(Motel ∞ ; school buses).

③ Two binary \wedge & \vee one unary complement

Same operations in propositional logic.

→ Boolean algebra

Venn
Dice
 2^n